

# Asynchronous games on Petri nets

Federica Adobbati

DISCo, Università degli studi di Milano-Bicocca

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# Concurrent systems

- A **user** interacts with an **environment**
- The user knows the structure of the system
- The user may not observe everything
- The user has a goal (perform an action, avoid a state, ...)
- The environment is hostile or indifferent

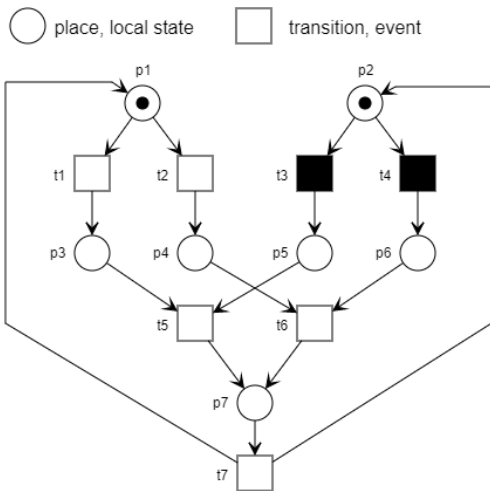
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## Applications

System control, model-checking, ...

## Modelling language: Petri nets

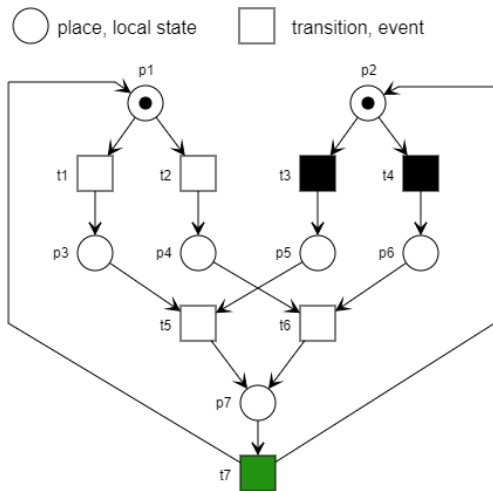


1-safe distributed nets

□ uncontrollable

■ controllable

## Modelling language: Petri nets



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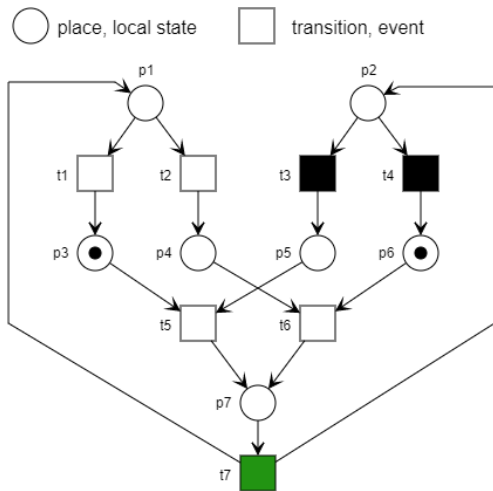
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## Example

Can the user force the occurrence of transition  $t_7$ , by controlling only black transitions?

## Modelling language: Petri nets



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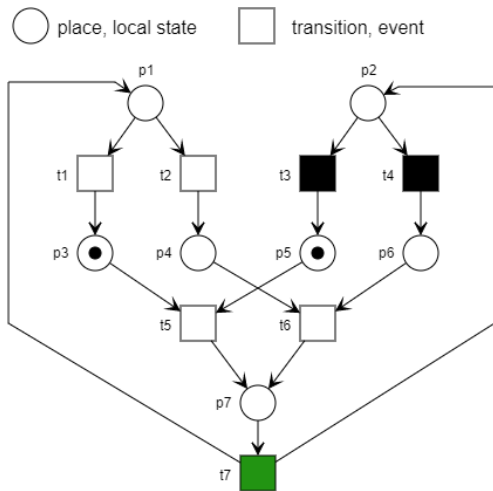
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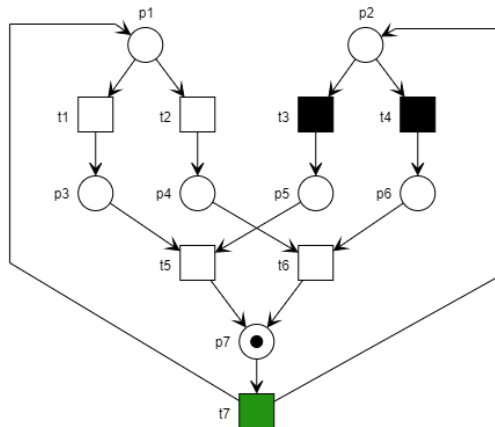
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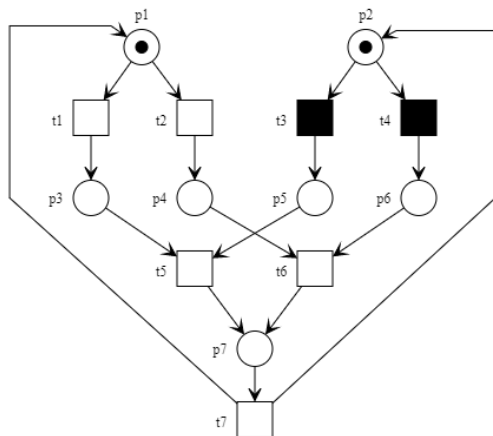
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# Asynchronous game on Petri nets



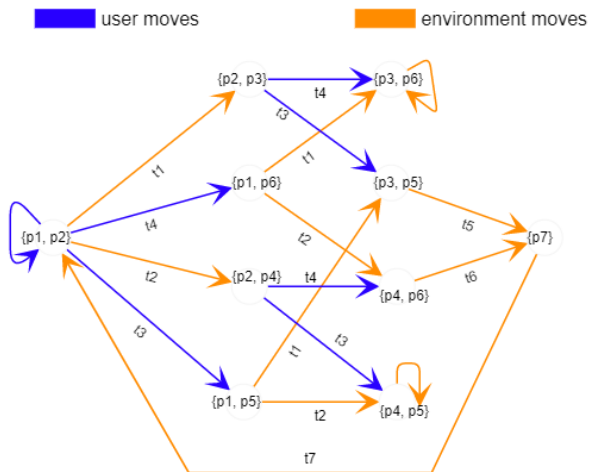
## Rules

- Whenever a transition is enabled, its owner can decide to 'fire' it
- The Environment must guarantee the progress of the system

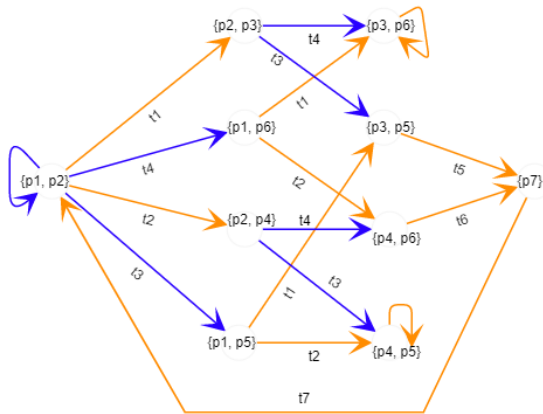
## Strategy

$$\alpha : \text{Mark}(N) \rightarrow 2^K$$

## Asynchronous game on concurrent game structure



## ATL (Alur, Henzinger, Kupferman 2002)



## Syntax

An ATL formula is one of the following:

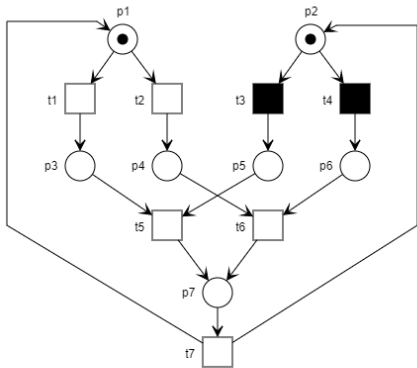
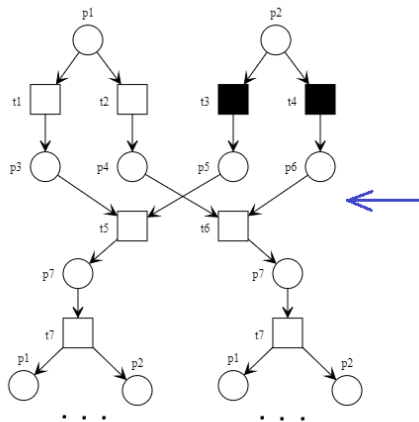
- a proposition  $p$
- $\neg\phi$  or  $\phi_1 \vee \phi_2$
- $\langle\langle A \rangle\rangle X\phi$ ,  $\langle\langle A \rangle\rangle G\phi$ ,  $\langle\langle A \rangle\rangle F\phi$ , or  $\langle\langle A \rangle\rangle\phi_1 U\phi_2$ , with  $A$  set of players

$\phi, \phi_1, \phi_2$  ATL formulas

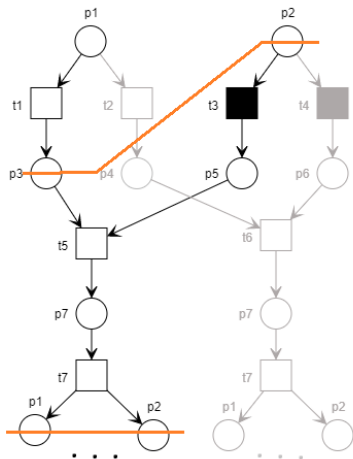
## Example

Is  $\langle\langle user \rangle\rangle Fp_7$  satisfied on this system?

# Game on the unfolding



# Game on the unfolding



- Run: a possible history of what happened in the net
- B-cut: a maximal set of pairwise concurrent places



# Algorithm for the reachability of a target transition

## Goal

The user reaches a target transition

## General idea

- Recursive algorithm on the unfolding, with backtracking
- Generation of a finite prefix of the unfolding
- In every marking, the strategy selects all the controllable transitions that are 'useful' to reach the target

## Mechanism of the algorithm

### Unfolding\_Exploration

From B-cut  $\gamma$  add events until:

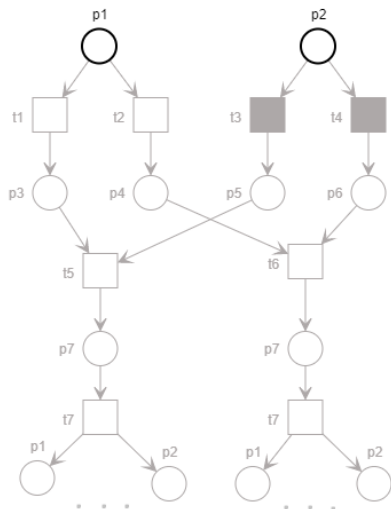
- Deadlock reached (the user tries to backtrack)
- Target fires (the environment tries to backtrack)
- Cycle detected (either no winning strategy or go to Explore\_Cycle)

### Explore\_Cycle

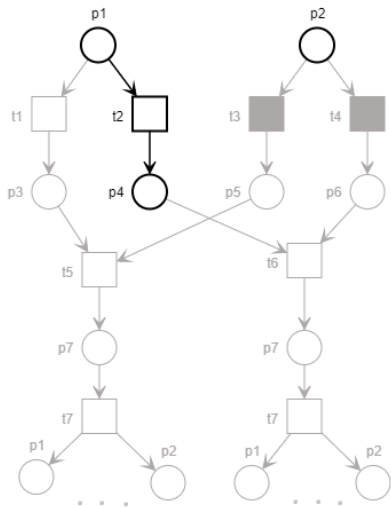
- Add a transition concurrent with the cycle
- Restart with Unfolding\_Exploration



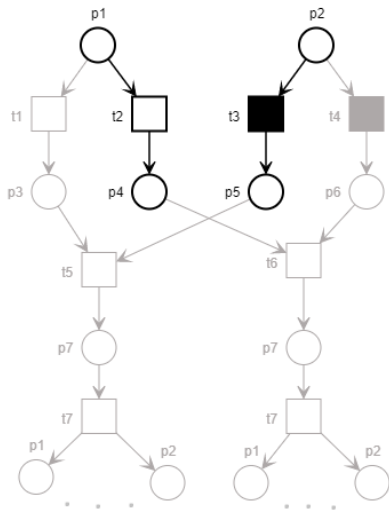
## Simulation of an execution



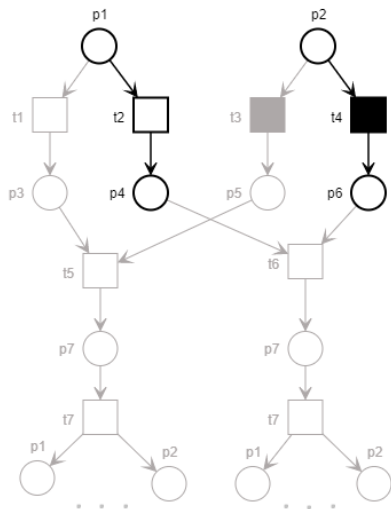
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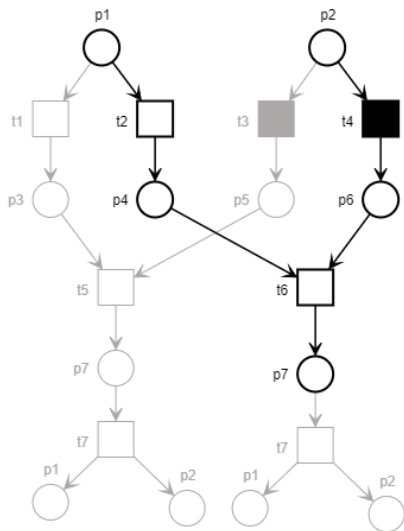
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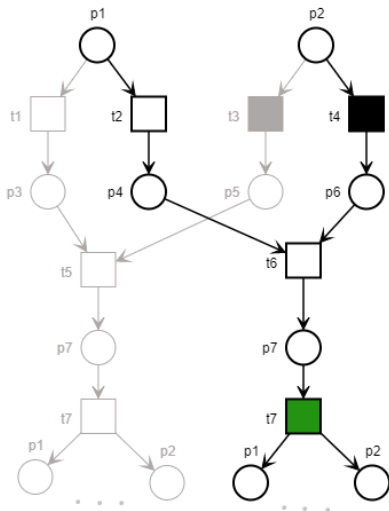
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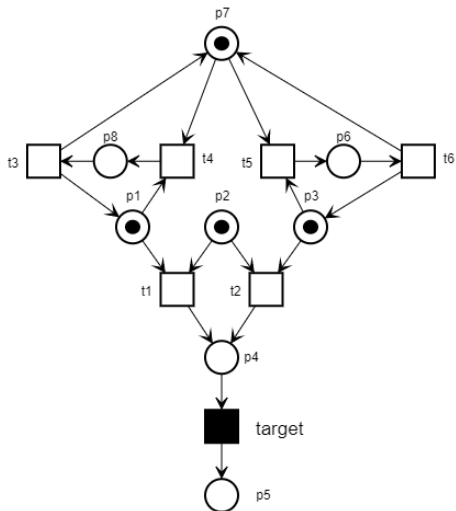


## Simulation of an execution



# Critical example

Solved with backtracking

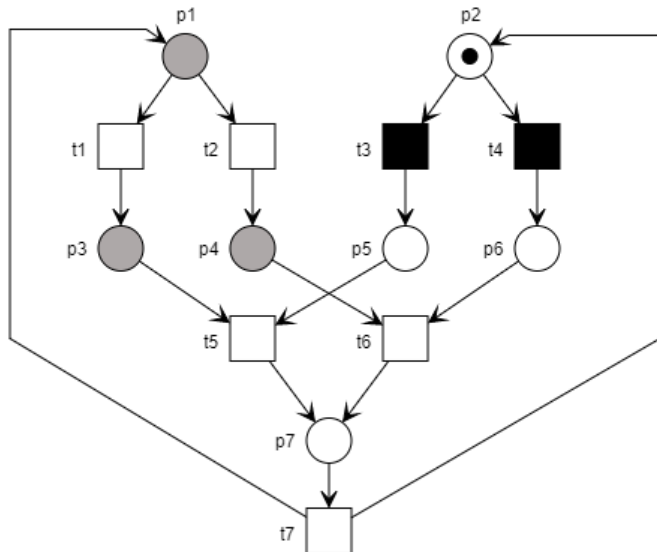


## Ongoing works and future developments

- Correctness and complexity of the algorithm
- Partial observability
- Implementability of the strategy



## Partial observability



# Implementable strategies

